### Lazy matrices for contraction-based algorithms

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Future work



The storage problem Lazy matrices Lazy matrices in quantum chemistry Future work 0000000 0000 000 000 000 000

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- 1 The storage problem
  - Problems with conventional approaches
  - contraction-based algorithms
- 2 Lazy matrices
  - The lazyten lazy matrix library
- 3 Lazy matrices in quantum chemistry
  - molsturm program package
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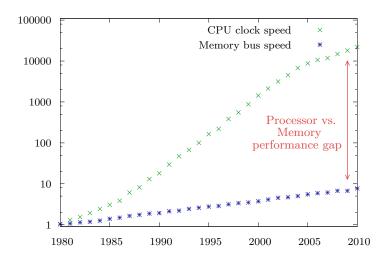
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# Processor vs. memory performance improvement



Future work

The storage problem

# Hartree-Fock equations

• Hartree-Fock equations

$$\left(-\frac{1}{2}\Delta + \hat{\mathcal{V}}_{\text{Nuc}} + \hat{\mathcal{V}}_{\text{H}}\left[\left\{\psi_{f}\right\}_{f \in I}\right] + \hat{\mathcal{V}}_{\text{x}}\left[\left\{\psi_{f}\right\}_{f \in I}\right] - \varepsilon_{f}\right)\psi_{f}(\underline{r}) = 0$$

with

$$-\frac{1}{2}\Delta$$
 Kinetic energy of electrons  $\hat{\mathcal{V}}_{\text{Nuc}}$  Electron-nuclear interaction

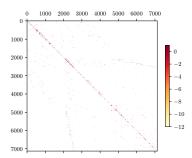
$$\hat{\mathcal{V}}_{\mathrm{H}}[\{\psi_f\}_{f\in I}]$$
 Hartree potential

$$\hat{\mathcal{V}}_{\mathbf{x}}[\{\psi_f\}_{f\in I}]$$
 Exchange potential

- Non-linear system of partial differential equations
- Non-linear eigenproblem for eigenpairs  $\{(\varepsilon_f, \psi_f)\}_{f \in I}$

#### Finite-element discretisation

- Finite elements: Piecewise polynomials with support only on a few neighbouring cells
- $\Rightarrow$  Need many finite elements (> 10<sup>6</sup>)
  - Typically sparse matrix structures:



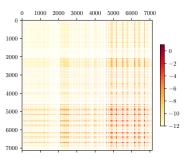
Typical discretisation of  $-\frac{1}{2}\Delta + \hat{\mathcal{V}}_{\text{Nuc}} + \hat{\mathcal{V}}_{\text{H}}$ 

#### Finite-element discretisation

#### Caveat

The storage problem

• But ...  $\hat{\mathcal{V}}_{\mathbf{x}}$  is so-called *non-local*:



**K**: Discretisation of  $\hat{\mathcal{V}}_{\mathbf{x}}$ 

- Building K takes  $\Omega(N^2)$  time and  $\mathcal{O}(N^2)$  storage
- Typically  $10^6 \cdot 10^6$  elements  $\approx 8 \, \text{TB}$  storage

#### Finite-element discretisation

contraction-based scheme

- Iterative solvers only need matrix-vector products
- Matrix-vector product of **K**: Theoretically  $\mathcal{O}(N)$
- ⇒ contraction-based or matrix-free algorithm:
  - Never build **K** in storage
  - Use expression for **K** to directly form contraction of matrix to vectors

### contraction-based SCF

#### Possible problems

• For SCF we will need

$$\mathbf{F} = \mathbf{T} + \mathbf{V}_{\mathrm{Nuc}} + \mathbf{J} + \mathbf{K} + \cdots$$

- Requirements differ:
  - Optimal storage scheme
  - Optimal contraction scheme
  - Approximations / costs
- ⇒ contraction expression complicated to code
- $\Rightarrow$  We would like to stay flexible

# Characteristics of contraction-based algorithms

### Advantages

- Scaling (storage and time) reduced in examples to  $\mathcal{O}(N)$
- Parallelisation easier
  - ⇒ Less data management
- Hardware trends are in favour

#### Disadvantages

- Matrices more intuitive than contraction-functions
- More computations

# Characteristics of contraction-based algorithms

### Advantages

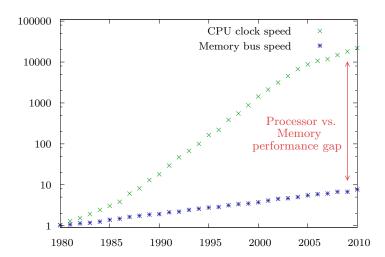
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#### Disadvantages

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- More computations
  - ⇒ Need efficient contraction schemes for the contraction
  - ⇒ Algorithms more complicated

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# Lazy matrices

- Stored matrix: All elements reside in memory
- Lazy matrix:
  - Generalisation of matrices
    - State
    - Non-linear
    - Elements may be expressions
  - ⇒ Obtaining elements expensive
    - Evaluation of internal expression: Delayed until contraction
    - For convenience: Offer matrix-like interface

# Using lazy matrices

• Program as usual

$$D = A + B$$

• Build expression tree internally

$$\boxed{ \mathbf{D} } = \boxed{ \mathbf{A} } + \boxed{ \mathbf{B} }$$

• On application:

$$\mathbf{D}\boldsymbol{x} = (\mathbf{A}\boldsymbol{x}) + (\mathbf{B}\boldsymbol{x})$$

# Using lazy matrices

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$$D = A + B$$

• Build expression tree internally

$$\boxed{ \mathbf{D} } = \boxed{ \mathbf{A} } + \boxed{ \mathbf{B} } = \boxed{ \mathbf{A}^{+}_{\mathbf{B}} }$$

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The storage problem

# Using lazy matrices

• Program as usual

$$D = A + B$$

• Build expression tree internally

$$\begin{bmatrix} \mathbf{D} \end{bmatrix} = \begin{bmatrix} \mathbf{A} \end{bmatrix} + \begin{bmatrix} \mathbf{B} \end{bmatrix} = \begin{bmatrix} \mathbf{A} \\ \mathbf{A} \end{bmatrix}$$

• On application:

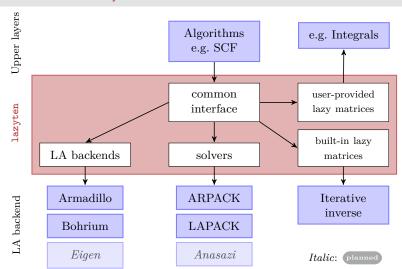
$$\mathbf{D}\boldsymbol{x} = (\mathbf{A}\boldsymbol{x}) + (\mathbf{B}\boldsymbol{x})$$

#### Notes and observations

- lazyten: Bookkeeping for contraction-functions
- Programmer still sees matrices
- ⇒ Language for writing contraction-based algorithms
  - Lazy matrices allow layered responsibility for computation, e.g.  $(\mathbf{A} + \mathbf{B})\boldsymbol{x}$ 
    - $\mathbf{A}\mathbf{x}$  and  $\mathbf{B}\mathbf{x}$  decided by implementation of  $\mathbf{A}$  and  $\mathbf{B}$
    - $(\mathbf{A}x) + (\mathbf{B}x)$  done in linear algebra backend
- ⇒ Proper modularisation between
  - Higher-level algorithms
  - Lazy matrix implementations
  - LA backends

### lazyten: Lazy matrix library

#### Structure of the library



A& O

• lazyten<sup>1</sup>: Prototype C++ implementation

```
1 typedef SmallVector < double > vector type;
2 typedef SmallMatrix < double > matrix type;
auto v = random < vector type > (100);
4 DiagonalMatrix < matrix type > diag(v);
5 auto a = random<matrix type>(100,100);
auto b = random<matrix type>(100,100);
7
8 // No computation: Just build expression tree
9 auto sum = diag + a;
auto projector = sum * inverse(sum);
auto tree = b - projector * b;
12
13 // Evaluate tree on application:
14 SmallVector < double > res = tree * v;
```

<sup>1</sup>https://lazyten.org

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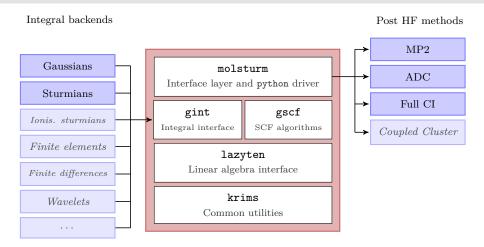


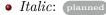


# Lazy matrices for quantum chemistry codes

- contraction-based algorithms
- ⇒ Lower memory footprint, scaling improvements
  - Readable code
- ⇒ Great for teaching or to play around
  - Abstraction between integrals and SCF algorithms
- ⇒ Plug and play integral libraries
- ⇒ Swap LA backends
- ⇒ Basis-type independent SCF / quantum chemistry

#### molsturm structure





# molsturm design

- Enables contraction-based SCF routines
- Flexiblity as primary goal
- Behaviour controlled via python
  - Keywords to change basis type or solver
  - All computed data available in numpy format
  - No input file, just a python script
- python utilities
  - Import / export results
  - Post-HF calculations

# molsturm interface: CCD residual (parts)

```
\begin{split} r_{ij}^{ab} &= -\frac{1}{2} \sum_{mnef} \left\langle mn||ef \right\rangle t_{mn}^{af} \, t_{ij}^{eb} + \frac{1}{2} \sum_{mnef} \left\langle mn||ef \right\rangle t_{mn}^{bf} \, t_{ij}^{ea} - \frac{1}{2} \sum_{mnef} \left\langle mn||ef \right\rangle t_{in}^{ef} \, t_{mj}^{ab} \\ &+ \frac{1}{2} \sum_{mnef} \left\langle mn||ef \right\rangle t_{jn}^{ef} \, t_{mi}^{ab} + \frac{1}{4} \sum_{mnef} \left\langle mn||ef \right\rangle t_{mn}^{ab} \, t_{ij}^{ef} + \frac{1}{2} \sum_{mnef} \left\langle mn||ef \right\rangle t_{im}^{ae} \, t_{jn}^{bf} \\ &- \frac{1}{2} \sum_{mnef} \left\langle mn||ef \right\rangle t_{jm}^{ae} \, t_{in}^{bf} - \frac{1}{2} \sum_{mnef} \left\langle mn||ef \right\rangle t_{im}^{be} \, t_{jn}^{af} + \frac{1}{2} \sum_{mnef} \left\langle mn||ef \right\rangle t_{jm}^{be} \, t_{in}^{af} \end{split}
```

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Outlook

The storage problem

# Lazy matrix expression optimisation

$$\begin{bmatrix} \mathbf{C} \\ + \\ \begin{pmatrix} \mathbf{A} \\ \end{pmatrix} + \begin{bmatrix} \mathbf{B} \\ \end{pmatrix} + \begin{bmatrix} \mathbf{A} \\ \mathbf{B} \end{bmatrix}$$

$$= \begin{bmatrix} \mathbf{A} \\ \mathbf{A} \\ \mathbf{B} \end{bmatrix}$$

- Matrix expression tree  $\equiv$  abstract syntax tree  $\equiv$  DAG
- ⇒ May be optimised by standard methods

Outlook

# Lazy matrix expression optimisation

$$\begin{bmatrix} \mathbf{C} \\ + \\ \begin{pmatrix} \mathbf{A} \\ \end{pmatrix} + \\ \begin{pmatrix} \mathbf{B} \\ \end{pmatrix} + \\ \begin{pmatrix} \mathbf{A} \\ \end{pmatrix} + \\ \begin{pmatrix} \mathbf{B} \\ \end{pmatrix} + \\ \begin{pmatrix} \mathbf{A} \\ \end{pmatrix} + \\ \begin{pmatrix} \mathbf{B} \\ \end{pmatrix} + \\ \begin{pmatrix} \mathbf{A} \\ \end{pmatrix} + \\ \begin{pmatrix} \mathbf{B} \\ \end{pmatrix} + \\ \begin{pmatrix} \mathbf{A} \\ \end{pmatrix} + \\ \begin{pmatrix} \mathbf{B} \\ \end{pmatrix} + \\ \begin{pmatrix} \mathbf{A} \\ \end{pmatrix} + \\ \begin{pmatrix} \mathbf{B} \\ \end{pmatrix} + \\ \begin{pmatrix} \mathbf{A} \\ \end{pmatrix} + \\ \begin{pmatrix} \mathbf{B} \\ \end{pmatrix} + \\ \begin{pmatrix} \mathbf{A} \\ \end{pmatrix} + \\ \begin{pmatrix} \mathbf{B} \\ \end{pmatrix} + \\ \begin{pmatrix} \mathbf{A} \\ \end{pmatrix} + \\ \begin{pmatrix} \mathbf{B} \\ \end{pmatrix} + \\ \begin{pmatrix} \mathbf{A} \\ \end{pmatrix} + \\ \begin{pmatrix} \mathbf{B} \\ \end{pmatrix} + \\ \begin{pmatrix} \mathbf{A} \\ \end{pmatrix} + \\ \begin{pmatrix} \mathbf{B} \\ \end{pmatrix} + \\ \begin{pmatrix} \mathbf{A} \\ \end{pmatrix} + \\ \begin{pmatrix} \mathbf{B} \\ \end{pmatrix} + \\ \begin{pmatrix} \mathbf{A} \\ \end{pmatrix} + \\ \begin{pmatrix} \mathbf{B} \\ \end{pmatrix} + \\ \begin{pmatrix} \mathbf{A} \\ \end{pmatrix} + \\ \begin{pmatrix} \mathbf{B} \\ \end{pmatrix} + \\ \begin{pmatrix} \mathbf{A} \\ \end{pmatrix} + \\ \begin{pmatrix} \mathbf{B} \\ \end{pmatrix} + \\ \begin{pmatrix} \mathbf{A} \\ \end{pmatrix} + \\ \begin{pmatrix} \mathbf{B} \\ \end{pmatrix} + \\ \begin{pmatrix} \mathbf{A} \\ \end{pmatrix} + \\ \begin{pmatrix} \mathbf{B} \\ \end{pmatrix} + \\ \begin{pmatrix} \mathbf{A} \\ \end{pmatrix} + \\ \begin{pmatrix} \mathbf{B} \\ \end{pmatrix} + \\ \begin{pmatrix} \mathbf{A} \\ \end{pmatrix} + \\ \begin{pmatrix} \mathbf{B} \\ \end{pmatrix} + \\ \begin{pmatrix} \mathbf{A} \\ \end{pmatrix} + \\ \begin{pmatrix} \mathbf{B} \\ \end{pmatrix} + \\ \begin{pmatrix} \mathbf{A} \\ \end{pmatrix} + \\ \begin{pmatrix} \mathbf{B} \\ \end{pmatrix} + \\ \begin{pmatrix} \mathbf{A} \\ \end{pmatrix} + \\ \begin{pmatrix} \mathbf{B} \\ \end{pmatrix} + \\ \begin{pmatrix} \mathbf{A} \\ \end{pmatrix} + \\ \begin{pmatrix} \mathbf{B} \\ \end{pmatrix} + \\ \begin{pmatrix} \mathbf{A} \\ \end{pmatrix} + \\ \begin{pmatrix} \mathbf{B} \\ \end{pmatrix} + \\ \begin{pmatrix} \mathbf{A} \\ \end{pmatrix} + \\ \begin{pmatrix} \mathbf{B} \\ \end{pmatrix} + \\ \begin{pmatrix} \mathbf{A} \\ \end{pmatrix} + \\ \begin{pmatrix} \mathbf{A}$$

- Matrix expression tree  $\equiv$  abstract syntax tree  $\equiv$  DAG
- ⇒ May be optimised by standard methods

### Selection of LA backend

- Right now: LA backend is compiled in
- e.g. Bohrium backend
  - Uses just-in-time (JIT) compilation
  - Very specific for hardware
  - Compilation takes time
- Better: Dynamic selection
  - Load on the expression tree
  - Availablility of backends
  - Hardware specs

# Extension to lazy tensors

- contraction is a special tensor contraction
- $\Rightarrow$  Lazy tensors:
  - Delay all tensor contractions as long as possible

e.g. 
$$\tilde{k}_{bf} = \sum_{\substack{acdo \mu\nu}} \mathcal{C}_{ao} \ c^{\mu}_{ab} I_{\mu\nu} c^{\nu}_{cd} \ \mathcal{C}_{co} \mathcal{C}_{df}$$

- Compare possible contraction schemes by complexity
- Execute cheapest evaluation scheme
- May incorporate and exploit symmetry
- Determine optimal contraction sequence automatically

- Dr. James Avery
- Prof. Andreas Dreuw and the Dreuw group



- Prof. Guido Kanschat
- HGS Mathcomp



### Questions?

The storage problem

- EMail: michael.herbst@iwr.uni-heidelberg.de
- Website/blog: https://michael-herbst.com
- Projects: https://lazyten.org and https://molsturm.org



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